# An algorithm for computing the fuzzy transitive closure of a bipolar weighted digraph 

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## 1 Notation and terminology

We have a bipolar weighted digraph $D=\left(V, w_{+}, w_{-}\right)$, where:

- $V=\{1,2, \ldots, n\}$ is the vertex set
- each arc is an ordered pair $(s, t)$, with the arrow from vertex $s$ to vertex $t$
- $w_{+}(s, t)$ is the weight of the positive arc $(s, t)$, and $0 \leq w_{+}(s, t) \leq 1$
- $w_{-}(s, t)$ is the weight of the negative arc $(s, t)$, and $0 \leq w_{-}(s, t) \leq 1$
- if the positive/negative arc $(s, t)$ is absent from $D$, we set $w_{+}(s, t)=0$ or $w_{-}(s, t)=0$, respectively

The bipolar weighted digraph $D=\left(V, w_{+}, w_{-}\right)$is represented by its matrices $A=\left(a_{s t}\right)_{s, t=1}^{n}$ and $B=\left(b_{s t}\right)_{s, t=1}^{n}$ :

- $A$ and $B$ are of dimension $n \times n$;
- their entries are $a_{s t}=w_{+}(s, t)$ and $b_{s t}=w_{-}(s, t)$.


## 2 Fuzzy transitive closure

The fuzzy transitive closure of $D=(V, w)$ is the bipolar weighted digraph $D^{*}=\left(V, w_{+}^{*}, w_{-}^{*}\right)$, where (informally):

- $w_{+}^{*}(s, t)$ is the maximum weight of a minimal $(s, t)$-walk of positive sign
- $w_{-}^{*}(s, t)$ is (in absolute value) the maximum weight of a minimal $(s, t)$-walk of negative sign
- the sign of a walk is the product of signs of all arcs traversed by the walk
- the weight of a walk is (in absolute value) the minimum of the weights of the arcs of the walk
- a minimal $(s, t)$-walk of positive/negative sign is an $(s, t)$-walk not properly contained in an $(s, t)$-walk of the same sign


## 3 Input

Positive integer $n$ and matrices $A=\left(a_{s t}\right)_{s, t=1}^{n}$ and $B=\left(b_{s t}\right)_{s, t=1}^{n}$ as described above.

## 4 Output

At the end of the algorithm, matrices $A=\left(a_{s t}\right)_{s, t=1}^{n}$ and $B=\left(b_{s t}\right)_{s, t=1}^{n}$, represent the fuzzy transitive closure of $D$. That is, $a_{s t}=w_{+}^{*}(s, t)$ and $b_{s t}=w_{-}^{*}(s, t)$ for all $s, t=1,2, \ldots, n$.

## 5 Algorithm

## procedure $\operatorname{FuzzyTC}(n, A, B)$

begin
Comment: compute the fuzzy transitive closure.
for $u=1,2, \ldots, n$ do Comment: $u$ is the new allowable vertex on the walk.
for $i=1,2$ do Comment: to allow for two traversals of vertex $u$.
begin
Comment: $A^{\prime}=\left(a_{s t}^{\prime}\right)_{s, t=1}^{n}$ and $B^{\prime}=\left(b_{s t}^{\prime}\right)_{s, t=1}^{n}$ will store the new entries of matrices $A$ and $B$, respectively.

> for $s=1,2, \ldots, n$ do
> for $t=1,2, \ldots, n$ do
> $\quad$ begin
> $\quad a_{s t}^{\prime}:=\max \left\{a_{s t}, \min \left(a_{s u}, a_{u t}\right), \min \left(b_{s u}, b_{u t}\right)\right\}$
> $b_{s t}^{\prime}:=\max \left\{b_{s t}, \min \left(a_{s u}, b_{u t}\right), \min \left(b_{s u}, a_{u t}\right)\right\}$
> end
for $s=1,2, \ldots, n$ do $\quad$ Comment: Update $A$ and $B$.
for $t=1,2, \ldots, n$ do
begin
$a_{s t}:=a_{s t}^{\prime}$
$b_{s t}:=b_{s t}^{\prime}$
end
end
output $A$
output $B$

Comment: now $A$ contains the positive and $B$ the absolute values of the negative weights of the fuzzy transitive closure.
end

## 6 Modification

If the input digraph contains no parallel arcs, of which one is positive and one negative (that is, there is no $(s, t)$ with $w_{+}(s, t)>0$ and $\left.w_{-}(s, t)>0\right)$, then the input can be a single $\operatorname{matrix} M=\left(m_{s t}\right)_{s, t=1}^{n}$ with entries

$$
m_{s t}= \begin{cases}w_{+}(s, t) & \text { if } w_{+}(s, t)>0 \\ -w_{-}(s, t) & \text { if } w_{-}(s, t)>0 \\ 0 & \text { otherwise }\end{cases}
$$

Input matrices $A$ and $B$ for procedure FuzzyTC can then be computed using the algorithm below.
procedure MatricesAB $(n, M)$

## begin

Comment: copy all positive entries of $M$ into matrix $A=\left(a_{s t}\right)_{s, t=1}^{n}$ and all negative ones (their absolute values) into matrix $B=\left(b_{s t}\right)_{s, t=1}^{n}$.
for $s=1,2, \ldots, n$ do
for $t=1,2, \ldots, n$ do begin
if $m_{s t} \geq 0$ then
begin

$$
\begin{aligned}
& a_{s t}:=m_{s t} \\
& b_{s t}:=0
\end{aligned}
$$

end
else
begin
$a_{s t}:=0$
$b_{s t}:=-m_{s t}$
end
end
output $A$
output $B$
end

